

CS 523: Computer Graphics, Spring 2009

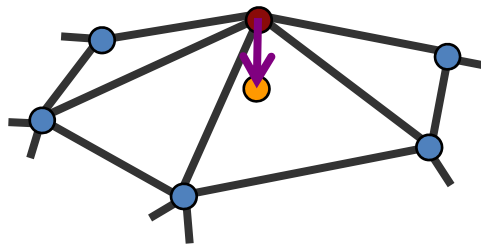
# Interactive Shape Modeling

Numerical methods

# Linear Solvers

## Motivation

- Laplace-type systems

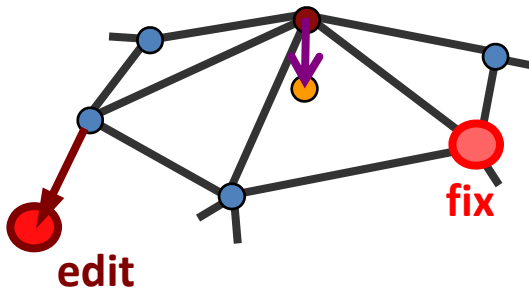


$$\delta_i = \sum_{j \in N(i)} w_{ij} (\mathbf{v}_i - \mathbf{v}_j)$$

$$\begin{array}{l} \mathbf{L} \mathbf{v}_x = \delta_x \\ \mathbf{L} \mathbf{v}_y = \delta_y \\ \mathbf{L} \mathbf{v}_z = \delta_z \end{array}$$

# Linear Solvers

## Motivation



$$\begin{array}{|c|} \hline L \\ \hline 1 \\ \hline 1 \\ \hline \end{array} \begin{array}{|c|} \hline v_x \\ \hline \end{array} = \begin{array}{|c|} \hline \delta_x \\ \hline c_x \\ \hline e_x \\ \hline \end{array}$$

$$\begin{array}{|c|} \hline L \\ \hline 1 \\ \hline 1 \\ \hline \end{array} \begin{array}{|c|} \hline v_y \\ \hline \end{array} = \begin{array}{|c|} \hline \delta_y \\ \hline c_y \\ \hline e_y \\ \hline \end{array}$$

$$\begin{array}{|c|} \hline L \\ \hline 1 \\ \hline 1 \\ \hline \end{array} \begin{array}{|c|} \hline v_z \\ \hline \end{array} = \begin{array}{|c|} \hline \delta_z \\ \hline c_z \\ \hline e_z \\ \hline \end{array}$$

# Linear Solvers

## Motivation

$$\begin{array}{c} \text{L} \\ 1 \\ 1 \end{array} \quad \mathbf{v}_x = \begin{array}{c} \delta_x \\ \mathbf{c}_x \\ \mathbf{e}_x \end{array}$$

$$\tilde{\mathbf{x}} = \arg \min_{\mathbf{x}} \left( \left\| \mathbf{L}\mathbf{x} - \boldsymbol{\delta}_x \right\|^2 + \sum_{s=1}^k |x_k - c_k|^2 \right)$$

... and the same for  $y$  and  $z$

# Linear Solvers

## Motivation

$$\begin{array}{|c|} \hline \mathbf{L} \\ \hline 1 \\ \hline 1 \\ \hline \end{array} \begin{array}{|c|} \hline \mathbf{v}_x \\ \hline \end{array} = \begin{array}{|c|} \hline \delta_x \\ \hline \mathbf{c}_x \\ \hline \mathbf{e}_x \\ \hline \end{array}$$

$$\tilde{\mathbf{L}} \mathbf{x} = \mathbf{c}$$

Normal Equations:

$$\begin{aligned} \tilde{\mathbf{L}}^T \tilde{\mathbf{L}} \mathbf{x} &= \tilde{\mathbf{L}}^T \mathbf{c} \\ \mathbf{x} &= (\tilde{\mathbf{L}}^T \tilde{\mathbf{L}})^{-1} \tilde{\mathbf{L}}^T \mathbf{c} \end{aligned}$$

# Linear Systems

- Matrix is often fixed, rhs changes

The diagram illustrates the equation  $Ax = b$ . On the left, a blue square labeled  $A$  is multiplied by a gray vertical rectangle labeled  $x$ . This is followed by an equals sign and another gray vertical rectangle labeled  $b$ . Below the  $A$  box, an upward-pointing arrow is labeled  $\tilde{L}^T \tilde{L}$ . Below the  $b$  box, an upward-pointing arrow is labeled  $\tilde{L}^T c$ .

# Iterative Solvers

- General approach: try to minimize some energy function  $E(\mathbf{x})$
- Linear case:  $E(\mathbf{x}) = \|\mathbf{Ax} - \mathbf{b}\|^2$
- Start from a guess  $\mathbf{x}_0$
- Iteratively improve:  $\mathbf{x}_{i+1} = g(\mathbf{x}_i)$
- Convergence:  $E(\mathbf{x})$  sufficiently small

# Descent Search

General algorithm

- Input: initial guess  $\mathbf{x}_0 \in \mathbb{R}^n$
- Step 0: set  $i = 0$
- Step 1: if  $E(\mathbf{x}) < \varepsilon$  stop,  
else compute *search direction*  $\mathbf{h}_i \in \mathbb{R}^n$
- Step 2: compute the *step size*  $\lambda_i$   
 $\lambda_i \in \arg \min_{\lambda \geq 0} E(\mathbf{x}_i + \lambda \cdot \mathbf{h}_i)$  ← Line search
- Step 3: set  $\mathbf{x}_{i+1} = \mathbf{x}_i + \lambda_i \mathbf{h}_i$ , goto Step 1



# Descent Search

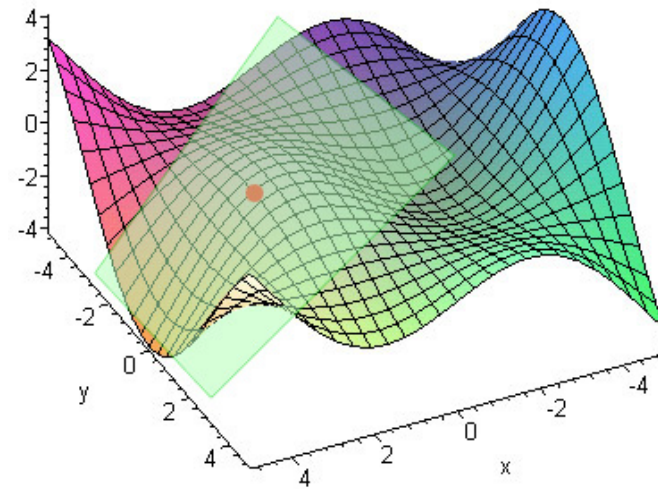
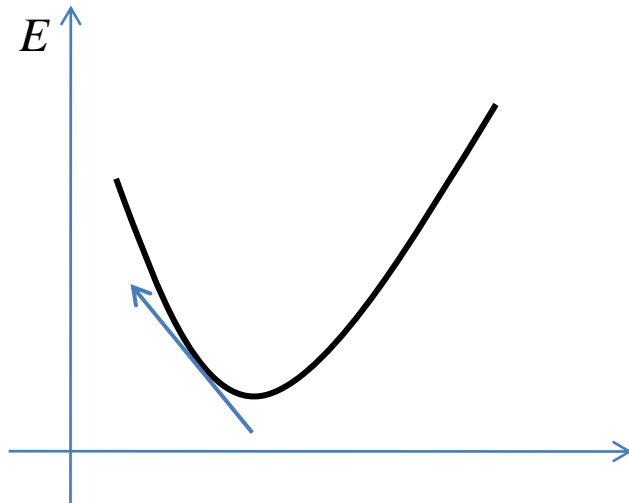
Quadratic energy (linear problem)

- Input: initial guess  $\mathbf{x}_0 \in \mathbb{R}^n$
- Step 0: set  $i = 0$
- Step 1: if  $\|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2 < \varepsilon$  stop,  
else compute *search direction*  $\mathbf{h}_i \in \mathbb{R}^n$
- Step 2: compute the *step size*  $\lambda_i$   
 $\lambda_i \in \arg \min_{\lambda \geq 0} \|\mathbf{A}(\mathbf{x}_i + \lambda \cdot \mathbf{h}_i) - \mathbf{b}\|$  ← Line search
- Step 3: set  $\mathbf{x}_{i+1} = \mathbf{x}_i + \lambda_i \mathbf{h}_i$ , goto Step 1

# Search Direction $\mathbf{h}_i$

Steepest descent

- Gradient is the direction in which the function grows the fastest

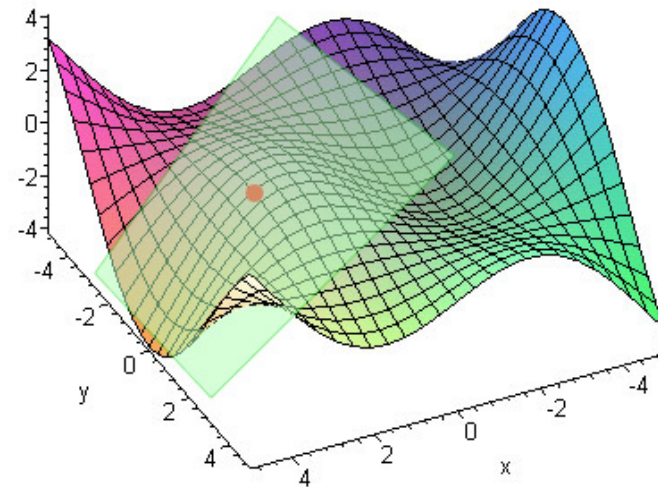
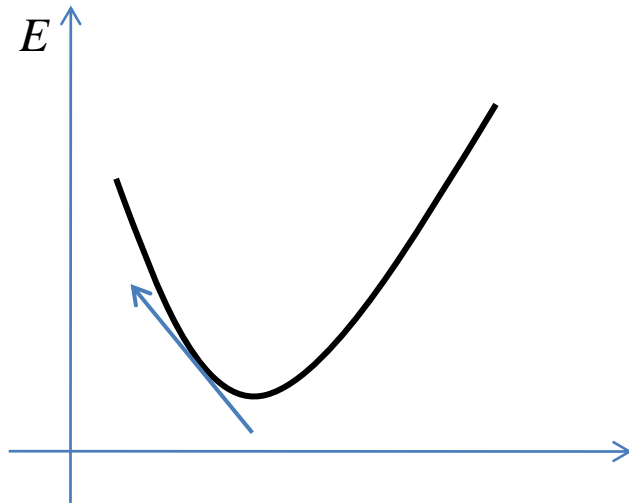


$$\mathbf{h}_i = -\nabla E(\mathbf{x}_i) / \|\nabla E(\mathbf{x}_i)\|$$

# Search Direction $\mathbf{h}_i$

Steepest descent

- Gradient is the direction in which the function grows the fastest

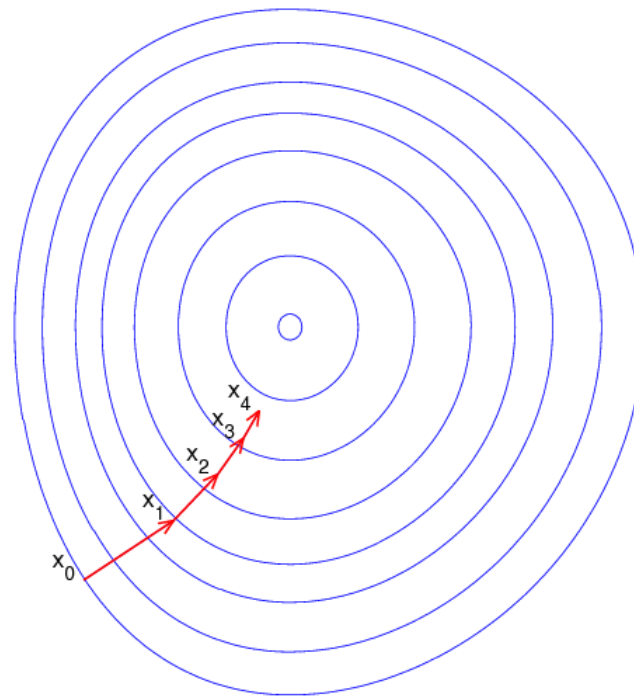


$$\nabla E(\mathbf{x}_i) = 2(\mathbf{A}^T \mathbf{A} \mathbf{x}_i - \mathbf{A}^T \mathbf{b})$$

# Search Direction $\mathbf{h}_i$

Steepest descent

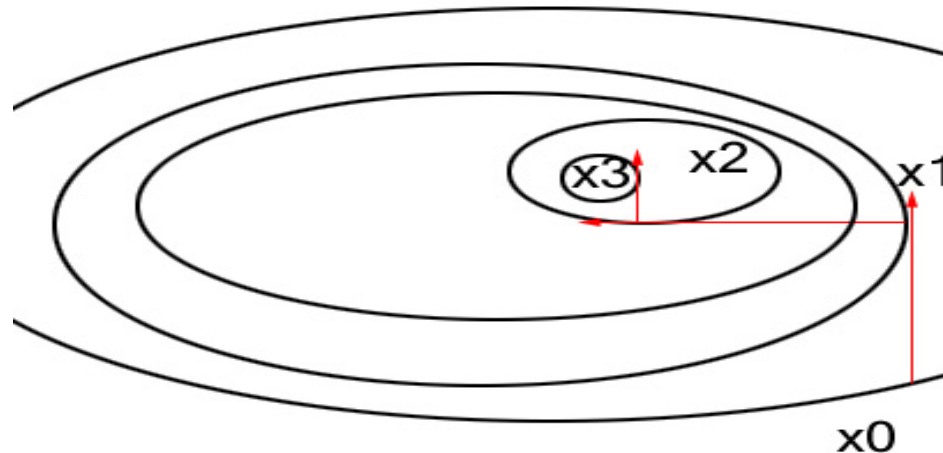
- Gradient is the direction in which the function grows the fastest



# Search Direction $\mathbf{h}_i$

Steepest descent

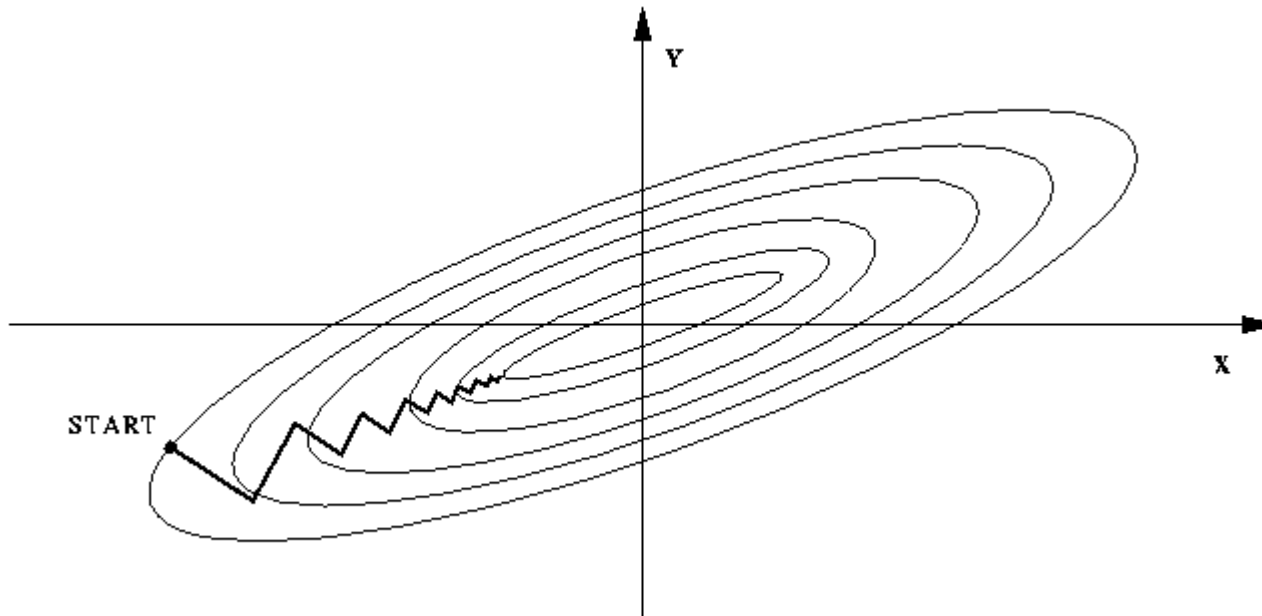
- Unlucky case: we pick the same direction many times



# Search Direction $\mathbf{h}_i$

Steepest descent

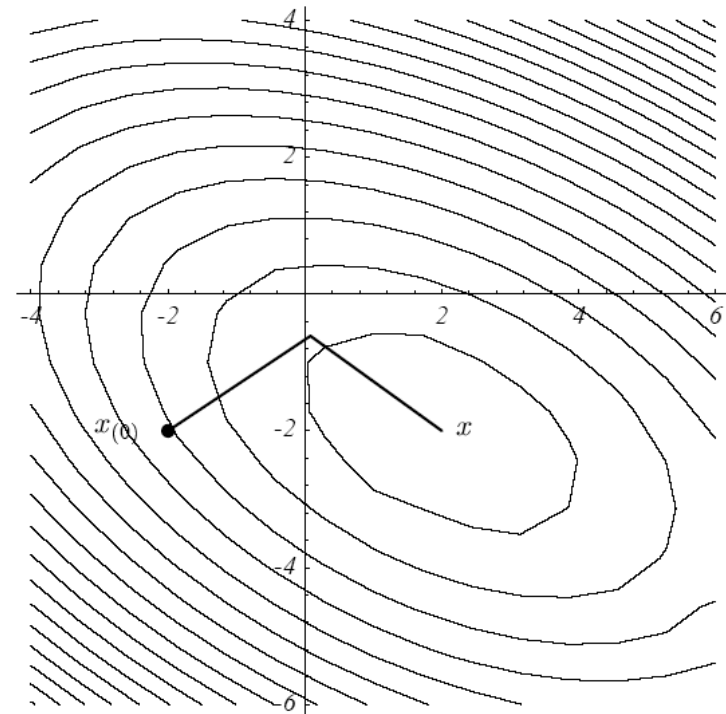
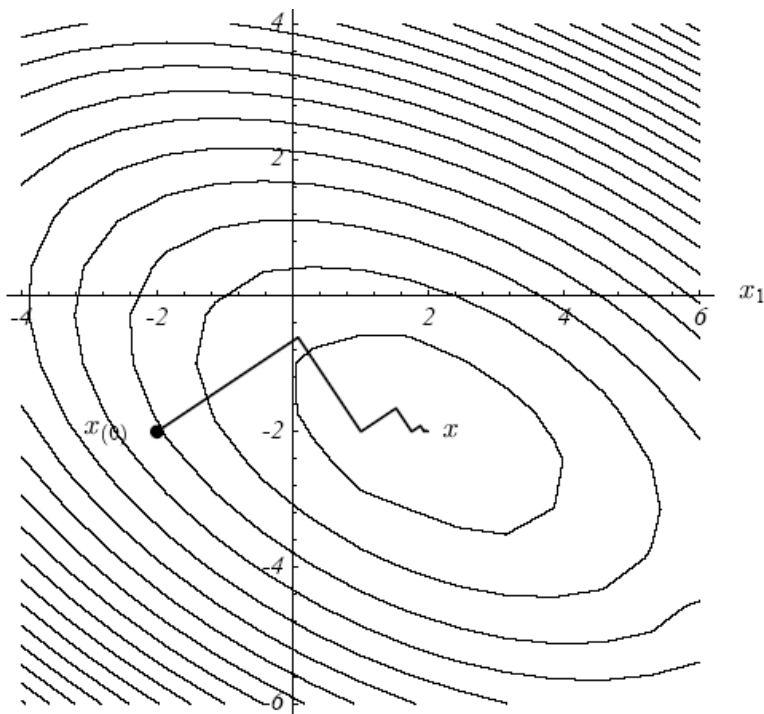
- Unlucky case: we pick the same direction many times



# Search Direction $\mathbf{h}_i$

Conjugate gradient

- Choose  $n$  linearly independent directions
- $\Rightarrow$  Converge in  $n$  steps



# Search Direction $\mathbf{h}_i$

Conjugate gradient

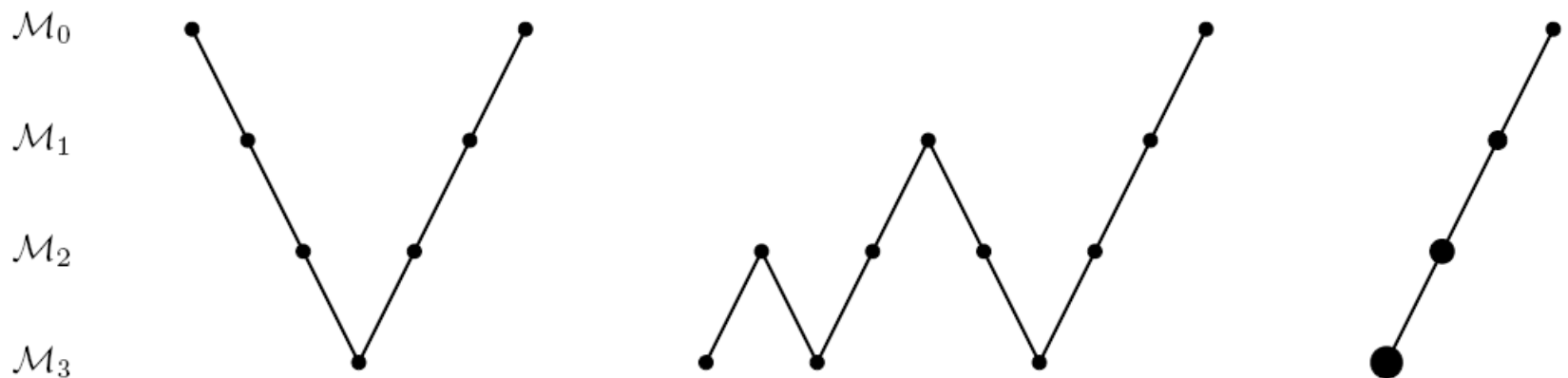
- The directions  $\mathbf{h}_1, \mathbf{h}_2, \dots, \mathbf{h}_n$  are chosen to be mutually “conjugate”, i.e., orthogonal w.r.t. the inner product defined by  $A$

$$\langle A\mathbf{h}_i, \mathbf{h}_j \rangle = \mathbf{h}_j^T A\mathbf{h}_i = 0$$



# Multigrid Solvers

- Coarsen the matrix and the rhs
- Solve on the coarse level, then interpolate to the finer level
- On meshes: geometric multigrid, i.e. coarsen the mesh by edge collapse operations



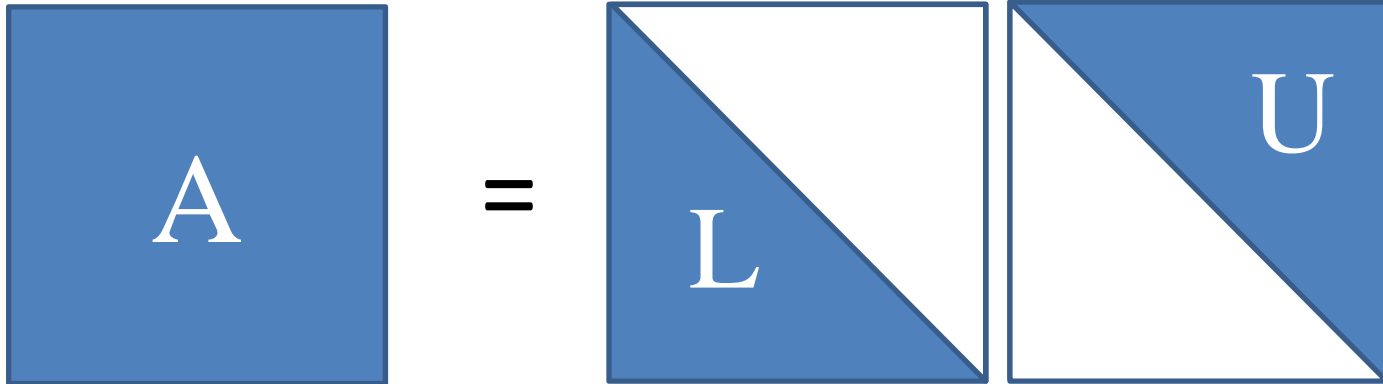
# Iterative Solvers

## Discussion

- Efficient in memory
  - Only store the matrix  $A$
- Not much gain when the rhs changes
  - Still need to iterate to find the solution, even though  $A$  is the same
- Too slow for interactive applications
- Problem-dependent parameters

# Matrix Factorization

LU decomposition

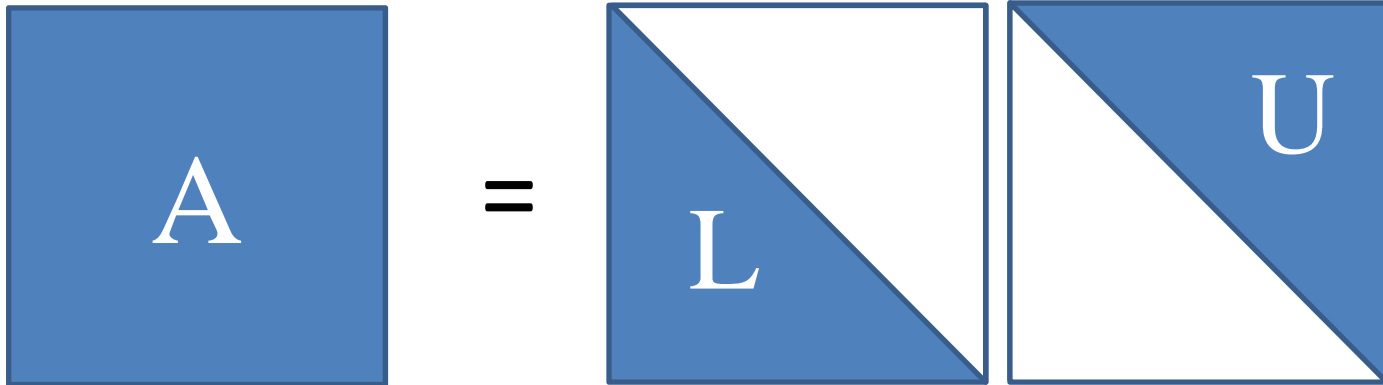


$$A\mathbf{x} = \mathbf{b}$$

$$LU\mathbf{x} = \mathbf{b}$$

# Matrix Factorization

LU decomposition

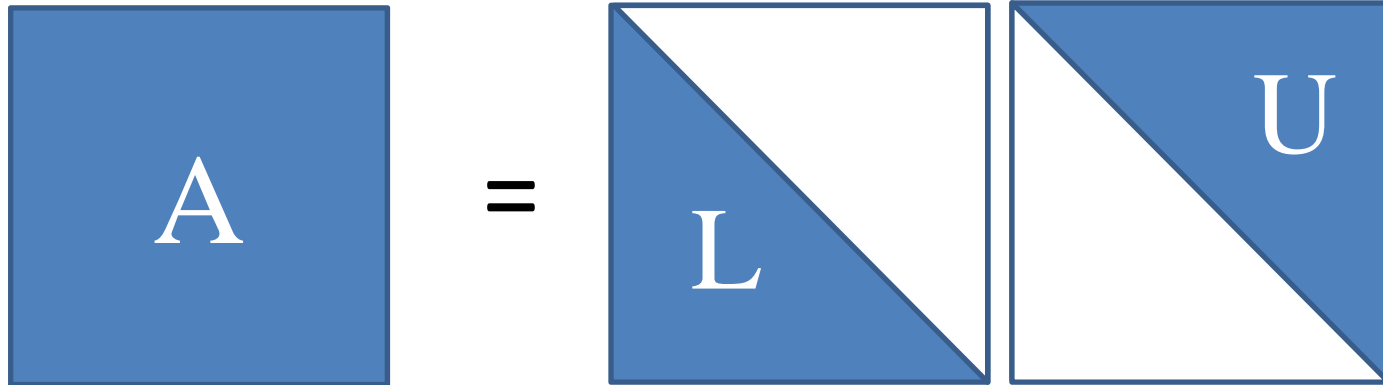


$$A\mathbf{x} = \mathbf{b}$$

$$L(U\mathbf{x}) = \mathbf{b}$$

# Matrix Factorization

LU decomposition



$$\begin{aligned} \mathbf{Ax} &= \mathbf{b} \\ \mathbf{L(Ux)} &= \mathbf{b} \end{aligned}$$



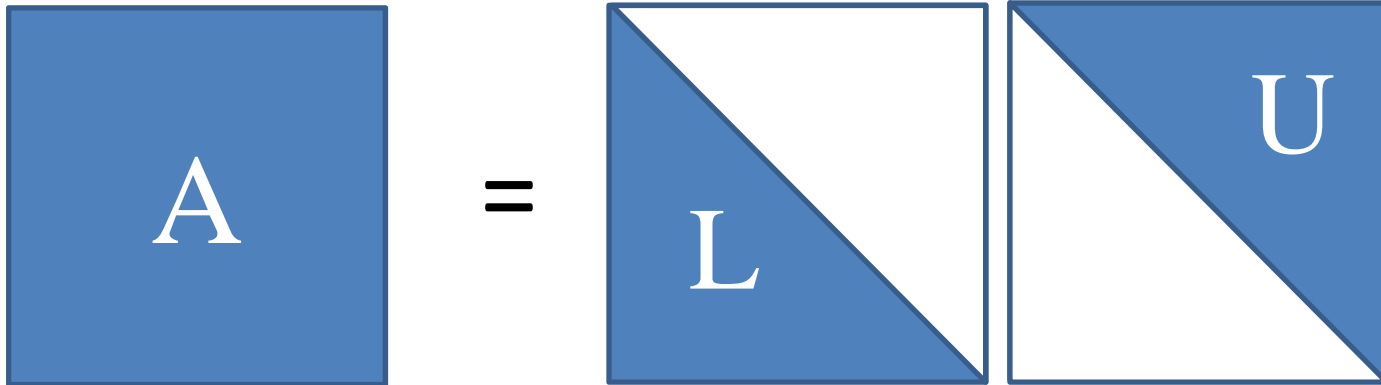
$$\begin{aligned} \mathbf{Ly} &= \mathbf{b} \\ \mathbf{Ux} &= \mathbf{y} \end{aligned}$$



This is backsubstitution.  
If  $L$ ,  $U$  are sparse it is very fast. The hard work is computing  $L$  and  $U$

# Matrix Factorization

LU decomposition



$$\mathbf{Ax} = \mathbf{b}$$
$$\mathbf{L(Ux)} = \mathbf{b}$$



$$\mathbf{y} = \mathbf{L}^{-1}\mathbf{b}$$
$$\mathbf{x} = \mathbf{U}^{-1}\mathbf{y}$$

This is backsubstitution.  
If  $\mathbf{L}$ ,  $\mathbf{U}$  are sparse it is very fast. The hard work is computing  $\mathbf{L}$  and  $\mathbf{U}$

# Matrix Factorization

## Cholesky decomposition

The diagram illustrates the Cholesky decomposition of a matrix  $A$ . On the left is a solid blue square labeled  $A$ . To its right is an equals sign. Further right are two square matrices. The first is a lower triangular matrix  $L$ , represented by a blue square with a white diagonal line from the top-left to the bottom-right. The second is the transpose of  $L$ ,  $L^T$ , represented by a white square with a blue diagonal line from the top-left to the bottom-right.

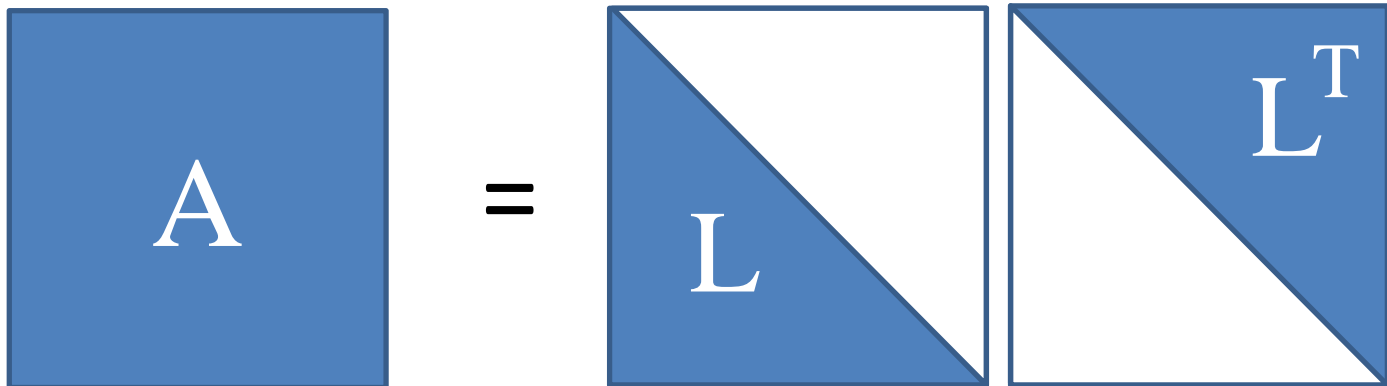
Cholesky factor exists if  $A$  is positive definite. It is even better than LU because we save memory.

# Cholesky Decomposition

$$A = LL^T$$

- $A$  is symmetric positive definite (PSD):

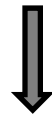
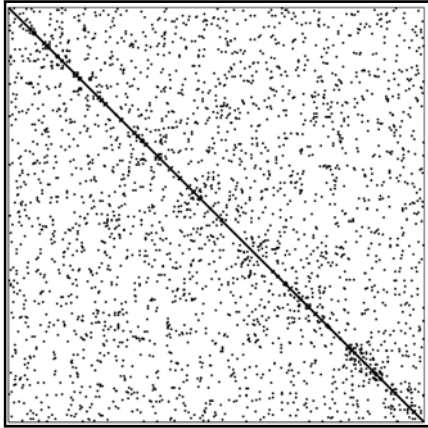
$$\forall \mathbf{x} \neq 0, \langle A\mathbf{x}, \mathbf{x} \rangle > 0 \quad \Leftrightarrow \quad \text{all } A\text{'s eigenvalues} > 0$$





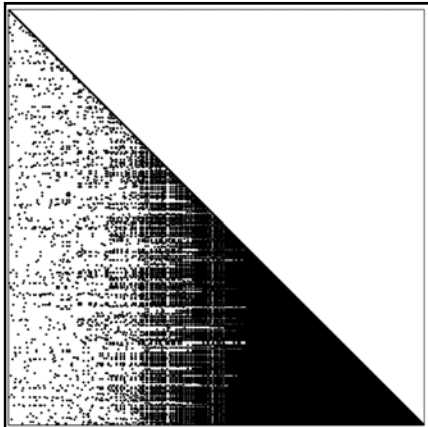
# Dense Cholesky Factorization

$A = LL^T$   
500×500 matrix  
3500 nonzeros



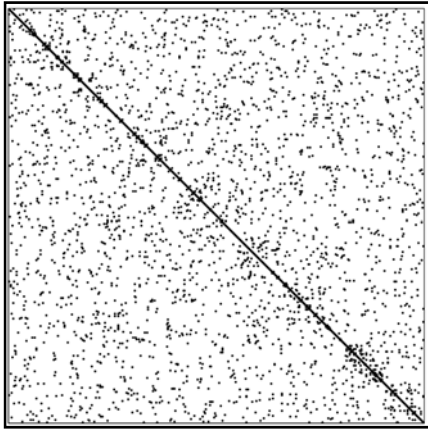
Cholesky Factorization

$L$   
36k nonzeros



# Sparse Cholesky Factorization

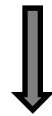
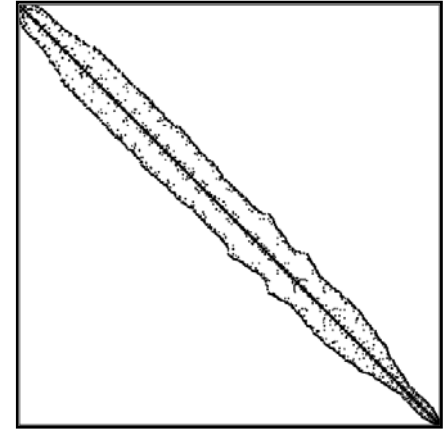
$A = LL^T$   
500×500 matrix  
3500 nonzeros



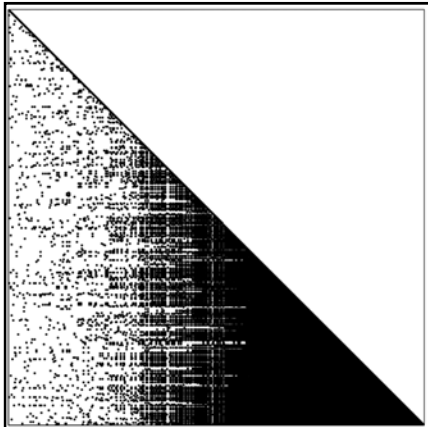
Reordering



$PAP^T$   
reverse Cuthill-  
McKee algorithm



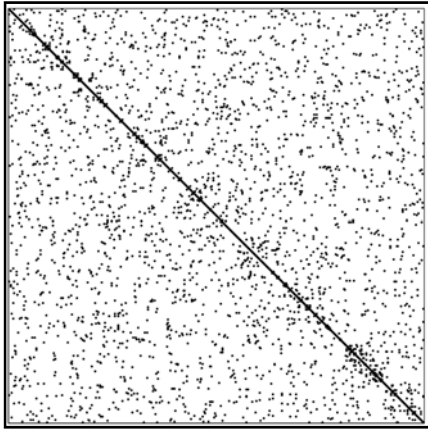
Cholesky Factorization



$L$   
36k nonzeros

# Sparse Cholesky Factorization

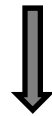
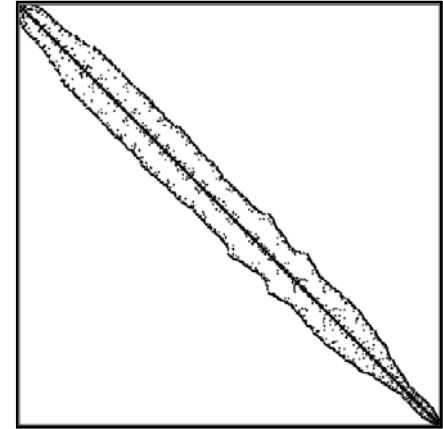
$A = LL^T$   
500×500 matrix  
3500 nonzeros



Reordering



$PAP^T$   
reverse Cuthill-  
McKee algorithm

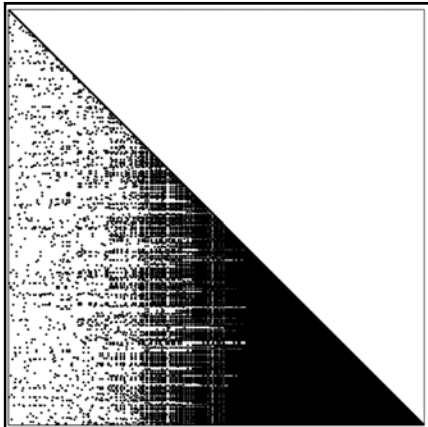


Cholesky Factorization



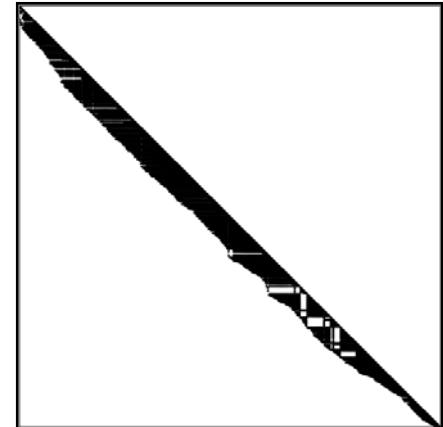
$L$

36k nonzeros



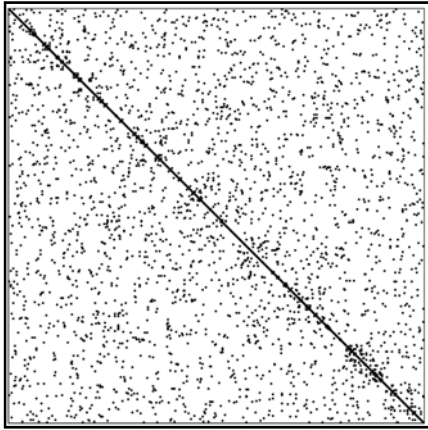
$L$

14k nonzeros



# Sparse Cholesky Factorization

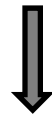
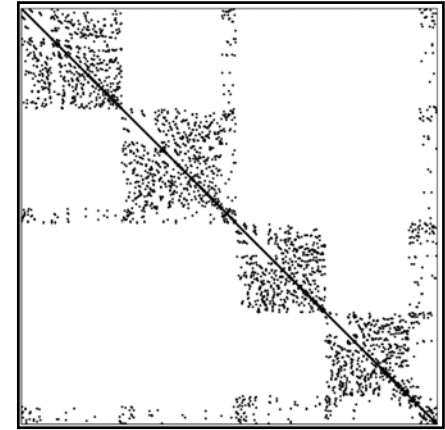
$A = LL^T$   
500×500 matrix  
3500 nonzeros



Reordering

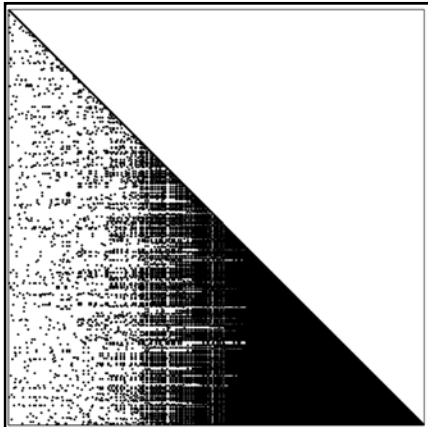


$PAP^T$   
nested dissection  
(parallelizable)



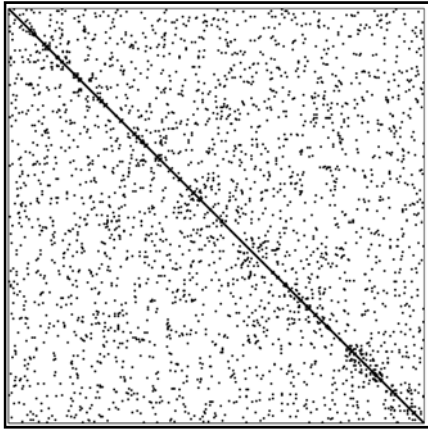
Cholesky Factorization

$L$   
36k nonzeros



# Sparse Cholesky Factorization

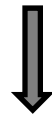
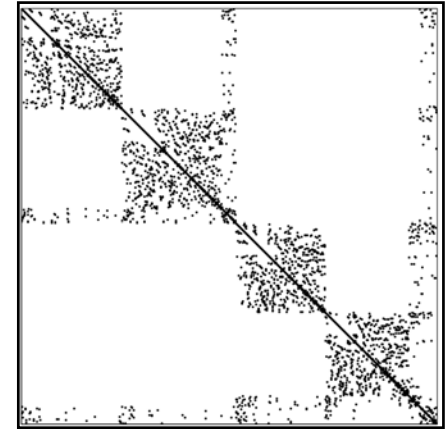
$A = LL^T$   
500x500 matrix  
3500 nonzeros



Reordering



$PAP^T$   
nested dissection  
(parallelizable)

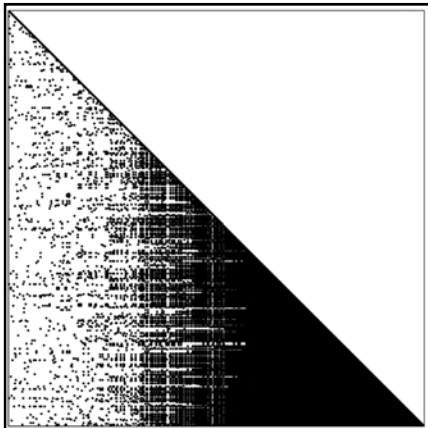


Cholesky Factorization



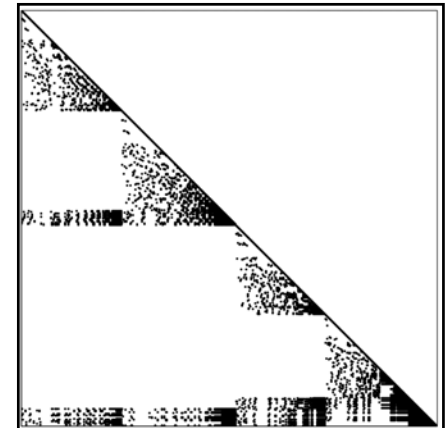
$L$

36k nonzeros



$L$

7k nonzeros



# Direct Solvers

## Discussion

- Highly accurate
  - Manipulate matrix structure
  - No iterations, everything is closed-form
- Easy to use
  - Off-the-shelf library, no parameters
- If  $A$  stays fixed, changing rhs ( $\mathbf{b}$ ) is cheap
  - Just need to back-substitute (factor precomputed)

# Direct Solvers

## Discussion

- High memory cost
  - Need to store the factor, which is typically denser than the matrix  $A$
- If the matrix  $A$  changes, need to re-compute the factor (expensive)

# TAUCS tutorial

- TAUCS: a library of sparse linear solvers
  - Has both iterative and direct solvers
  - Direct (Cholesky and LU) use reordering and are very fast
- I provide a wrapper for TAUCS on Sakai
  - Unsupported (= Windows, VS 2003)



# TAUCS tutorial

- Basic operations:
  - Define a sparse matrix structure
  - Fill the matrix with its nonzero values (i, j, v)
  - Factor  $A^T A$
  - Provide an rhs and solve

# TAUCS tutorial

- Basic operations:
  - Define a sparse matrix structure

```
InitTaucsInterface();  
  
int idA;  
idA = CreateMatrix(4, 3);
```

#rows #cols



# TAUCS tutorial

- Basic operations:
  - Fill the matrix A with its nonzero values (i, j, v)

```
SetMatrixEntry(idA, i, j, v);
```

# TAUCS tutorial

- Basic operations:
  - Fill the matrix A with its nonzero values (i, j, v)

```
SetMatrixEntry(idA, i, j, v);
```




matrix ID, obtained in CreateMatrix

# TAUCS tutorial

- Basic operations:
  - Fill the matrix A with its nonzero values (i, j, v)

```
SetMatrixEntry(idA, i, j, v);
```



row index i, column index j,  
zero-based

# TAUCS tutorial

- Basic operations:
  - Fill the matrix A with its nonzero values (i, j, v)

```
SetMatrixEntry(idA, i, j, v);
```

value of matrix entry ij  
for instance,  $-w_{ij}$

# TAUCS tutorial

- Basic operations:
  - Factor the matrix  $A^T A$

```
FactorATA( idA ) ;
```

# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

```
taucsType b[4] = {3, 4, 5, 6};  
taucsType x[3];  
  
SolveATA(idA, b, x, 1);
```



# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

```
taucsType b[4] = {3, 4, 5, 6};  
taucsType x[3];  
  
SolveATA(idA, b, x, 1);
```

↑  
typedef for double

# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

```
taucsType b[4] = {3, 4, 5, 6};  
taucsType x[3];  
  
SolveATA(idA, b, x, 1);
```

ID of the A matrix

# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

```
taucsType b[4] = {3, 4, 5, 6};  
taucsType x[3];  
  
SolveATA(idA, b, x, 1);
```

rhs for the LS system  $Ax = b$

# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

```
taucsType b[4] = {3, 4, 5, 6};  
taucsType x[3];  
  
SolveATA(idA, b, x, 1);
```

array for the solution

# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

A is 4x3

```
taucsType b[4] = {3, 4, 5, 6};  
taucsType x[3];  
  
SolveATA(idA, b, x, 1);
```

number of rhs's

# TAUCS tutorial

- Basic operations:
  - Provide an rhs and solve

A is 4x3

```
taucsType b2[8] = {3, 4, 5, 6, 7, 8, 9, 10};  
taucsType xy[6];  
  
SolveATA(idA, b2, xy, 2);
```

number of rhs's

# TAUCS tutorial

- If the matrix  $A$  is square a priori, no need to solve the LS system
- Then just use `FactorA()` and `SolveA()`

# Further Reading

- **Efficient Linear System Solvers for Mesh Processing**

Mario Botsch, David Bommes, Leif Kobbelt  
Invited paper at IMA Mathematics of Surfaces XI, Lecture  
Notes in Computer Science, Vol 3604, 2005, pp. 62-83.